

# Error correction codes



EduClash  
Just Another Way To Learn

# Block Error Correction Codes

- Error detection requires blocks to be retransmitted.
- This is inadequate for wireless communication for two reasons.
- The bit error rate on a wireless link can be quite high, which would result in a large number of retransmissions.
- In some cases, especially satellite links, the propagation delay is very long compared to the transmission time of a single frame.
- The result is a very inefficient system.
- It is desirable to be able to correct errors without requiring retransmission. Using the bits that were transmitted.

# Hamming code

- The number of redundant bits can be calculated using the following formula:
- $2^r > m + r + 1$  where,  $r$  = redundant bit,  $m$  = data bit
- Suppose the number of data bits is 7, then the number of redundant bits can be calculated using:
  - $2^4 > 7 + 4 + 1$
  - Thus, the number of redundant bits= 4

- **Example**

Hamming code=(8,4)

8 bit data block=00111001

(a) Transmitted block

|                   |      |      |      |      |      |      |      |      |      |      |      |      |
|-------------------|------|------|------|------|------|------|------|------|------|------|------|------|
| Bit Position      | 12   | 11   | 10   | 9    | 8    | 7    | 6    | 5    | 4    | 3    | 2    | 1    |
| Position Number   | 1100 | 1011 | 1010 | 1001 | 1000 | 0111 | 0110 | 0101 | 0100 | 0011 | 0010 | 0001 |
| Data Bit          | D8   | D7   | D6   | D5   |      | D4   | D3   | D2   |      | D1   |      |      |
| Check Bit         |      |      |      |      | C8   |      |      |      | C4   |      | C2   | C1   |
| Transmitted Block | 0    | 0    | 1    | 1    | 0    | 1    | 0    | 0    | 1    | 1    | 1    | 1    |
| Codes             |      |      | 1010 | 1001 |      | 0111 |      |      |      | 0011 |      |      |

(b) Check bit calculation prior to transmission

| Position          | Code |
|-------------------|------|
| 10                | 1010 |
| 9                 | 1001 |
| 7                 | 0111 |
| 3                 | 0011 |
| XOR = C8 C4 C2 C1 | 0111 |

(c) Received block

|                 |      |      |      |      |      |      |      |      |      |      |      |      |
|-----------------|------|------|------|------|------|------|------|------|------|------|------|------|
| Bit Position    | 12   | 11   | 10   | 9    | 8    | 7    | 6    | 5    | 4    | 3    | 2    | 1    |
| Position Number | 1100 | 1011 | 1010 | 1001 | 1000 | 0111 | 0110 | 0101 | 0100 | 0011 | 0010 | 0001 |
| Data Bit        | D8   | D7   | D6   | D5   |      | D4   | D3   | D2   |      | D1   |      |      |
| Check Bit       |      |      |      |      | C8   |      |      |      | C4   |      | C2   | C1   |
| Received Block  | 0    | 0    | 1    | 1    | 0    | 1    | 1    | 0    | 1    | 1    | 1    | 1    |
| Codes           |      |      | 1010 | 1001 |      | 0111 | 0110 |      |      | 0011 |      |      |

(d) Check bit calculation after reception

| Position       | Code |
|----------------|------|
| Hamming        | 0111 |
| 10             | 1010 |
| 9              | 1001 |
| 7              | 0111 |
| 6              | 0110 |
| 3              | 0011 |
| XOR = syndrome | 0110 |

# BCH

- Bose, Chaudhuri, and Hocquenghem (BCH) codes are among the most powerful cyclic block codes and are widely used in wireless applications.
- For any positive pair of integers  $m$  and  $t$ , there is a binary  $(n, k)$  BCH code with the following parameters:
  - Block length:  $n = 2^m - 1$
  - Number of check bits:  $n - k \leq mt$
  - Minimum distance:  $d_{\min} \geq 2t + 1$
- This code can correct all combinations of  $t$  or fewer errors.

# Reed solomon codes

- Reed-Solomon (RS) codes are a widely used subclass of nonbinary BCH codes.
- With RS codes, data are processed in chunks of  $m$  bits, called symbols.
- An  $(n, k)$  RScode has the following parameters:
  - Symbol length:  $m$  bits per symbol
  - Block length:  $n = 2^m - 1$  symbols
  - Data length:  $k$  symbols
  - Size of check code:  $n - k = 2t$  symbols =  $m(2t)$  symbols
  - Minimum distance:  $d_{\min} = 2t + 1$  symbols

Let  $t = 1$  and  $m = 2$ . Denoting the symbols as 0, 1, 2, 3 we can write their binary equivalents as  $0 = 00$ ;  $1 = 01$ ;  $2 = 10$ ;  $3 = 11$ . The code has the following parameters.

$$n = 2^2 - 1 = 3 \text{ symbols} = 6 \text{ bits}$$

$$(n - k) = 2 \text{ symbols} = 4 \text{ bits}$$

This code can correct any burst error that spans a symbol of 2 bits.

# Application and example

- Storage devices (including tape, Compact Disk, DVD, barcodes, etc)
  - Wireless or mobile communications (including cellular telephones, microwave links, etc)
  - Satellite communications
  - Digital television / DVB
  - High-speed modems such as ADSL, xDSL, etc.
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- A popular Reed-Solomon code is RS(255,223) with 8-bit symbols. Each codeword contains 255 code word bytes, of which 223 bytes are data and 32 bytes are parity. For this code:
    - $n = 255, k = 223, s = 8$
    - $2t = 32, t = 16$
    - The decoder can correct any 16 symbol errors in the code word: i.e. errors in up to 16 bytes anywhere in the codeword can be automatically corrected.
    - Given a symbol size  $s$ , the maximum codeword length ( $n$ ) for a Reed-Solomon code is  $n = 2s - 1$
    - For example, the maximum length of a code with 8-bit symbols ( $s=8$ ) is 255 bytes.



# Atmospheric Absorption

- At frequencies above 10 GHz, radio waves propagating through the atmosphere are subject to molecular absorption.
- The absorption as a function of frequency is very uneven
- There is a peak of water vapor absorption at around 22 GHz and a peak of oxygen absorption near 60 GHz.
- Vegetation and rain too effect the waves.