

# Queues

Rear of  
Queue

Front of  
Queue



# Queues

- A linear list of elements in which
  - Inserting can take place at the other end called **rear**.
  - deletion can take place only at one end called **front**.
- Queue principle : FIFO
- Conditions
  - If there is no element in queue/ empty queue
    - **FRONT = -1 ; REAR = -1**



# Queues

– If there is exactly one element in queue

- $\text{FRONT} = \text{REAR}$



- Whenever an element is added to the queue,  
 $\text{REAR} = \text{REAR} + 1$  or  $\text{REAR}++$
- If the item being added is the first element,  
 $\text{FRONT} = 0$ , indicating the queue is no longer empty.

# Operations on queues

- Four basic operations on queue

- **Enqueue :**

- Queue insert operation is called as **enqueue**.
    - The data is inserted in the rear.
    - If there is not room to insert another data in the queue , the queue is said to be in an **overflow state**.
    - Diagram
    - algorithm



Algorithm Enqueue(struct  
queue \*q, int item)

To add the data to the queue  
using array implementation

pre: struct stack \*s : pointer  
to the queue structure

item : data to be pushed  
in the stack

post : to add data to the queue

- 1.[check if array is full]  
if (q->rear=ARR-1)  
print "\n QUEUE IS FULL"
2. [increment rear and add item  
to the array]
  1. q->rear++
  2. q->a[q->rear]=item;
  3. count++;
3. [if the array is empty  
(front =-1) ,set front=0]  
if(q->front =-1)  
q->front=0

# Operations on queues

## – Dequeue:

- Queue delete operation is called as **dequeue**.
- Data is removed from the front of the queue.
- If there are no data in the queue, when a dequeue is attempted , then the queue is in an **underflow state**.
- Diagram
- algorithm



Algorithm int dequeue(struct queue \*q)

To remove data from the queue using array implementation

pre: struct stack \*s : pointer to the queue structure

post: data is removed from the queue and returned to the calling program

1. [if the queue is empty return null]  
if(q->front == -1)  
return NULL
2. [fetch data from the front of the queue]
  1. data = q->a[q->front]
  2. q->a[q->front] = NULL;
3. [ if last element was removed from the queue,  
if(q->front == q->rear)  
q->front = q->rear = -1;  
else  
q->front++;  
return data;



# Operations on queues

## – QueueFront:

- Returns data which is at the front of the queue without changing the contents of the queue.
- If there are no data in the queue, when a queue front is attempted , then the queue is in an **underflow state**.
- Diagram
- Algorithm





Algorithm int queuefront(struct  
queue \*q)

Assignment

To display the data which is in the  
front of the queue

pre: struct stack \*s : pointer to the  
queue structure

Post : data which is in the front of  
the queue is returned to the  
calling program



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# Operations on queues

## – QueueRear:

- Returns data which is at the rear of the queue without changing the contents of the queue.
- If there are no data in the queue, when a queue rear is attempted , then the queue is in an **underflow state**.
- Diagram
- Algorithm



## Assignment

Algorithm int queue rear(struct queue \*q)

To display the data which is in the rear end of the queue

pre: struct stack \*s : pointer to the queue structure

Post : data which is in the rear end of the queue is returned to the calling program



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# Display the contents of the queue

Algorithm displayqueue(struct queue \*q)

To display the contents of the queue

Pre : struct queue \*s : pointer to the queue structure

1.[initialize]

    x=0

2.[if rear=-1 then the queue is empty]

    1. if(q->rear==-1)

        print "QUEUE IS EMPTY"

        return;

    2. repeat while(x<=q->rear)

        1. printf q->a[x]

        2. x=x+1

# Circular Queues

- Disadvantage of linear queues:
  - If an item is removed from the linear array, the space remains unutilized. To overcome this, the circular array is used.
  - Algorithm:



# Algorithm to insert data in a Circular Queue

Algorithm ins\_circular\_q(Struct  
queue \*cq, int item)

To insert data into circular queue

Pre :

Post:

1. [Check if the queue is full]

If  $q \rightarrow \text{REAR} = \text{MAX} - 1$  and

$q \rightarrow \text{Front} = 0$

print "circular queue is  
full"

return

2. If  $q \rightarrow \text{REAR} = \text{MAX} - 1$

$q \rightarrow \text{REAR} = 0$

else

increment  $q \rightarrow \text{REAR}$

if( $q \rightarrow a[q \rightarrow \text{rear}] == 0$  ||

$q \rightarrow a[q \rightarrow \text{rear}] == -999$ )

$q \rightarrow a[q \rightarrow \text{REAR}] = \text{item}$

else

printf("QUEUE IS FULL");

3. [If an item is inserted , then  
queue is not empty]

if  $q \rightarrow \text{front} = -1$

$q \rightarrow \text{front} = 0$

## Algorithm to delete data from a Circular Queue

Algorithm int

delete\_from\_circular\_queue(

{struct queue \*q)

1. [Check if the queue is empty]

if(q->front == -1)

return

2. Assign the value of array to the variable data and set the array element to 0

data=a[q->front]

a[q->front]=0

3.[check if the queue is empty]

if(q->front==q->rear)

q->front=q->rear=-1

else

if(q->front==MAX-1)

q->front=0;

else

q->front++

return data

# Priority Queues

- Collection of elements where each element is assigned a priority and the order in which elements are deleted and processed.
- The following are the rules:
  - All elements of higher priority is before any element of lower priority.
  - 2 elements with the same priority are processed according to the order in which they are added to the queue.



# Priority Queues

## – Eg: Time Sharing System

- Programs of high priority are processed first and programs with the same priority form a standard queue.

- Eg:
- Algorithm



# Algorithm for inserting data in the priority queue

Algorithm add\_priority\_jobs(struct priorityqueue \*pq, struct data dt)

Post: to add jobs in a priority queue

1.[declare variables]

struct data temp

int i,j

2. if(pq->rear==MAX -1)

display " QUEUE IS FULL"

return

3. rear++;

d[pq->rear]=dt;

if pq->front== -1

    pq->front =0;

4. Repeat for(i=0;i<=pq->rear;i++)

1. Repeat for(j=i+1; j<= pq-> rear;j++)

    A. If(priority1 >priority2)

        swap the jobs

    B. If(priority1=priority2)

        1. if(orderno1>orderno2)

            swap the jobs



# Algorithm for displaying data in the priority queue

```
void displaycircularqueue(struct circularqueue *q)
```

```
1[initialize]
```

```
    x=0
```

```
2. Repeat while(x<=MAX -1)
```

```
    1. display q->a[x]
```

```
    2. x++
```

